

# On Double Inverse Semigroups

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# In the beginning. . .

In his 2006 paper “Note on commutativity in double semigroups and two-fold monoidal categories”, Kock introduced the notion of a double semigroup, along with some commutativity properties of them. In particular, he defines double inverse semigroups.

## Definition

A *double semigroup*  $(S, \odot, \odot)$  is a set equipped with two associative binary operations satisfying the *middle-four interchange law*: for all  $a, b, c, d \in S$ ,

$$(a \odot b) \odot (c \odot d) = (a \odot c) \odot (b \odot d).$$

- Horizontal product:  $a \odot b = \boxed{a \mid b}$ .

- Vertical product:  $a \odot b = \boxed{a \mid b}$ .

- Middle-four:

|     |     |
|-----|-----|
| $a$ | $b$ |
| $c$ | $d$ |

## Example

Any set  $D$  can be made into a double semigroup by equipping it with left and right projection:

$$a \odot b = a \text{ and } a \oslash b = b.$$

Associative:

$$\begin{array}{|c|c|c|} \hline a & b & c \\ \hline \end{array} = \begin{array}{|c|} \hline c \\ \hline \end{array} \quad \begin{array}{|c|} \hline a \\ \hline b \\ \hline c \\ \hline \end{array} = \begin{array}{|c|} \hline a \\ \hline \end{array}$$

Middle-four interchange law:

$$\begin{array}{|c|c|} \hline a & b \\ \hline c & d \\ \hline \end{array} = \begin{array}{|c|} \hline b \\ \hline \end{array}$$

## Definition

Given an element,  $x$  in a semigroup  $(S, \odot)$ ,  $x$  said to have an *inverse*  $x^\odot$  if

$$x = x \odot x^\odot \odot x \text{ and } x^\odot = x^\odot \odot x \odot x^\odot.$$

A semigroup is said to be an *inverse semigroup* if every element has a unique inverse. A double semigroup is said to be inverse if both of its operations are.

## Note

$x \odot x^\odot$  and  $x^\odot \odot x$  are idempotents:

- $(x \odot x^\odot) \odot (x \odot x^\odot) = (x \odot x^\odot \odot x) \odot x^\odot = x \odot x^\odot$
- $(x^\odot \odot x)(x^\odot \odot x) = (x^\odot \odot x \odot x^\odot) \odot x = x^\odot \odot x$

## Theorem (Kock)

*Double inverse semigroups are commutative.*

In the comments of his  $\text{\LaTeX}$  source code, Kock mentions that he does not have any “significant” examples of a double inverse semigroup. We aim to either find one, or to characterise double inverse semigroups.

## Coming soon:

- Explore Lawson's correspondence between inductive groupoids and inverse semigroups given by a pair of constructions.
- Define double inductive groupoids.
- Extend these constructions to double inductive groupoids and double inverse semigroups and establish an analogous correspondence.

## A quick notational note:

If  $f : A \rightarrow B$  is an arrow in a category:

### Notation

- *Domain of  $f$  :  $f \text{ dom} = A$ .*
- *Codomain of  $f$  :  $f \text{ cod} = B$ .*
- *Denote the composite*

$$A \xrightarrow{f} B \xrightarrow{g} C$$

*as  $f; g$  or  $fg$ .*

## Definition

Let  $(G, \bullet)$  be a groupoid and let  $\leq$  be a partial order defined on the arrows of  $G$ . We call  $(G, \bullet, \leq)$  and *ordered groupoid* whenever the following conditions are satisfied:

- If  $x \leq y$ , then  $x^{-1} \leq y^{-1}$ .
- If  $x \leq y$ ,  $u \leq v$ , then  $xu \leq yv$ .

## Note

*Identification of identity arrows with objects:*

- Gives  $\leq$  on objects

## Definition (cont'd)

- Let  $f \in G_1$  and let  $e$  be an object in  $G$  such that  $e \leq f \text{ dom}$ . Then there is a unique element  $(e_*|f) \in G_1$ , called the restriction of  $f$  by  $e$ , such that  $(e_*|f) \leq f$  and  $(e_*|f) \text{ dom} = e$ .
- Let  $f \in G_1$  and let  $e$  be an object in  $G$  such that  $e \leq f \text{ cod}$ . Then there is a unique element  $(f|_*e) \in G_1$ , called the corestriction of  $f$  by  $e$ , such that  $(f|_*e) \leq f$  and  $(f|_*e) \text{ cod} = e$ .

$$f \text{ dom} \xrightarrow{f} f \text{ cod}$$

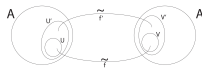
IV

$$e \xrightarrow{(e_*|f)} (e_*|f) \text{ cod}$$

## Example

Let  $A$  be a set. Construct an inductive groupoid with the following data:

- Objects :  $\mathcal{P}A$
- Arrows: Partial isomorphisms  $f : U \xrightarrow{\sim} V$  between subsets  $U, V \in \mathcal{P}A$
- $(f : U \rightarrow V) \leq (f' : U' \rightarrow V')$  if and only if  $U \subseteq U'$  and  $f'$  restricted to  $U$  (as functions) is  $f$ .



## Definition

An ordered groupoid  $G$  is an *inductive groupoid* if its objects form a meet-semilattice.

# Inductive Groupoids from Inverse Semigroups

## Construction

*Given an inverse semigroup  $(S, \odot)$  with the natural partial ordering  $\leq$ , define an inductive groupoid,  $(IG(S), \bullet)$ , with the following data:*

**Objects:** *idempotents of  $S$ ;  $IG(S)_0 = E(S)$ .*

**Arrows:** *elements of  $S$ .*

## Construction (cont'd)

**Arrows:** *elements of  $S$ .*

- $\text{sdom} = s \odot s^\odot$
- $\text{scod} = s^\odot \odot s$
- If  $a^\odot \odot a = b \odot b^\odot$ , define  $a \bullet b = a \odot b$
- Every arrow is an isomorphism with  $a^{-1} = a^\odot$
- $(a|_* e) = a \odot e$
- $(e_*| a) = e \odot a$

# Inverse Semigroups from Inductive Groupoids

## Construction

*Given an inductive groupoid  $(G, \bullet, \leq, \wedge)$ , construct an inverse semigroup  $(IS(G), \odot)$  with  $IS(G) = G_1$  and, for any  $a, b \in S$ ,*

$$a \odot b = (a|_* a \text{cod} \wedge b \text{dom}) \bullet (a \text{cod} \wedge b \text{dom}_* | b).$$

# An Isomorphism of Categories

## Notation

*Denote the category of inverse semigroups and semigroup homomorphisms as **IS**. Denote the category of inductive groupoids and inductive functors as **IG**.*

## Theorem (Lawson)

*The categories **IG** and **IS** are isomorphic.*

# GOAL: Double this theorem

## Definition

A *double category*  $\mathcal{D}$  consists of the following data:

- A collection  $\mathcal{D}_0$  of objects.
- A collection  $\text{Ver}(\mathcal{D})$  of vertical arrows.

Associative and unitary composition:

$$A \xrightarrow{\bullet f} B \xrightarrow{\bullet g} C = A \xrightarrow{\bullet f \bullet g} C$$

$$A \xrightarrow{\bullet 1_A} A \xrightarrow{\bullet f} B = A \xrightarrow{\bullet f} B = A \xrightarrow{\bullet f} B \xrightarrow{\bullet 1_B} B$$

## Definition (cont'd)

- A collection  $\text{Hor}(\mathcal{D})$  of horizontal arrows.

Associative and unitary composition:

$$A \xrightarrow{f} B \xrightarrow{g} C = A \xrightarrow{f \circ g} C$$

$$A \xrightarrow{\text{id}_A} A \xrightarrow{f} B = A \xrightarrow{f} B = A \xrightarrow{f} B \xrightarrow{\text{id}_B} B$$

## Definition (Cont'd)

- A collection  $\text{DbI}(\mathcal{D})$  of double cells. A double cell  $\alpha$  has the following form:

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow u & \alpha & \downarrow v \\ C & \xrightarrow{g} & D \end{array}$$

- $A, B, C$  and  $D$  are objects of  $\mathcal{D}$ .
- Horizontal domain and codomain:

$$\alpha \text{hdom} = u \text{ and } \alpha \text{hcod} = v$$

- Vertical domain and codomain:

$$\alpha \text{vdom} = f \text{ and } \alpha \text{vcod} = g$$

## Definition (cont'd)

These double cells must come together with:

- An associative and unitary horizontal composition,  $\circ$ .
- An associative and unitary vertical composition,  $\bullet$ .
- Horizontal and vertical composition of double cells must satisfy the middle-four interchange law. That is, for any  $\alpha, \beta, \gamma, \delta \in \text{DbI}(\mathcal{D})$ ,

$$(\alpha \bullet \beta) \circ (\gamma \bullet \delta) = (\alpha \circ \gamma) \bullet (\beta \circ \delta).$$

## Definition

A *double inductive groupoid*, denoted DIG,

$$\mathcal{G} = (\text{Obj}(\mathcal{G}), \text{Ver}(\mathcal{G}), \text{Hor}(\mathcal{G}), \text{Dbl}(\mathcal{G}))$$

is a double groupoid (i.e., a double category in which every vertical arrow, horizontal arrow and double cell is an isomorphism) such that :

## Definition (cont'd)

$(\text{Ver}(\mathcal{G}), \text{Dbl}(\mathcal{G}))$  is an inductive groupoid.

- Composition: horizontal composition,  $\circ$ .
- Partial order :  $\leq$ .
- Meet of vertical arrows  $e$  and  $f$  :  $e \wedge_h f$ .
- For a double cell  $\alpha$  and a vertical arrow  $e$  with  $e \leq \alpha \text{hdom}$ , horizontal restriction :  $(e_* | \alpha)$ .
- If  $e \leq \alpha \text{hcod}$ , horizontal corestriction:  $(\alpha |_* e)$ .

## Definition (cont'd)

$(\text{Hor}(\mathcal{G}), \text{Dbl}(\mathcal{G}))$  is an inductive groupoid.

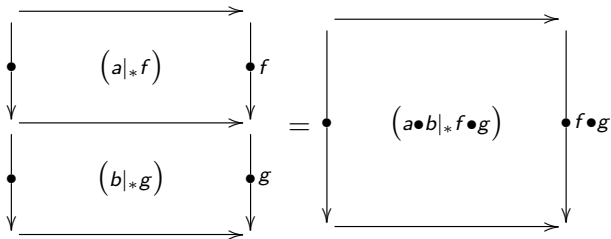
- Composition: vertical composition,  $\bullet$ .
- Partial order :  $\lesssim$ .
- Meet of horizontal arrows  $e$  and  $f$  :  $e \wedge_v f$ .
- For a double cell  $\alpha$  and a horizontal arrow  $e$  with  $e \lesssim \alpha \text{vdom}$ , vertical restriction :  $[e_* | \alpha]$ .
- If  $e \lesssim \alpha \text{vcod}$ , vertical corestriction:  $[\alpha |_* e]$ .

## Definition (cont'd)

If  $a, b$  are double cells,  $f', g'$  are horizontal arrows and  $f, g$  are vertical arrows, the following laws about restrictions and corestrictions preserving composition hold:

- ①  $(a \bullet b|_* f \bullet g) = (a|_* f) \bullet (b|_* g).$
- ②  $[a \circ b|_* f' \circ g'] = [a|_* f'] \circ [b|_* g'].$
- ③  $(f \bullet g_*| a \bullet b) = (f_*| a) \bullet (g_*| b).$
- ④  $[f' \circ g'_*| a \circ b] = [f'_*| a] \circ [g'_*| b].$

$$(a \bullet b|_* f \bullet g) = (a|_* f) \bullet (b|_* g)$$

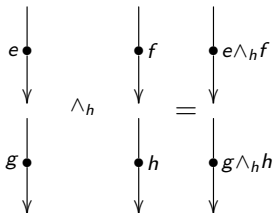


## Definition (cont'd)

If  $e, f, g$  and  $h$  are horizontal arrows and  $e', f', g'$  and  $h'$  are vertical arrows, the following laws about composition and meets satisfying middle-four hold:

- (a)  $(e \wedge_v f) \circ (g \wedge_v h) = (e \circ g) \wedge_v (f \circ h).$
- (b)  $(e' \wedge_h f') \bullet (g' \wedge_h h') = (e' \bullet g') \wedge_h (f' \bullet h').$

$$(e \wedge_h f) \bullet (g \wedge_h h) = (e \bullet g) \wedge_h (f \bullet h)$$



## Definition (cont'd)

If  $e$  and  $g$  are horizontal arrows  $f$  and  $h$  are objects, then the following rule about corestrictions and meets satisfying middle-four holds:

$$(e|_* f) \wedge_v (g|_* h) = (e \wedge_v g|_* f \wedge_v h)$$

$$(e|_*f) \wedge_v (g|_*h) = (e \wedge_v g|_*f \wedge_v h)$$

$$\begin{array}{ccc} \xrightarrow{(e|_*f)} & & f \\ & \wedge_v & \\ & = & \xrightarrow{(e \wedge_v g|_*f \wedge_v h)} f \wedge_v h \\ & & \\ \xrightarrow{(g|_*h)} & & h \end{array}$$

Similarly,

- (a)  $(e|_*f) \wedge_v (g|_*h) = (e \wedge_v g|_*f \wedge_v h).$
- (b)  $[e'|_*f'] \wedge_h [g'|_*h'] = [e' \wedge_h g'|_*f' \wedge_h h'].$
- (c)  $(e_*|f) \wedge_v (g_*|h) = (e \wedge_v g_*|f \wedge_v h).$
- (d)  $[e'_*|f'] \wedge_h [g'_*|h'] = [e' \wedge_h g'_*|f' \wedge_h h'].$

## Definition (cont'd)

If  $a$  is a double cell,  $f$  a horizontal arrow,  $g$  a vertical arrow and  $x$  an object such that

$$f \lesssim a \text{vcod}$$

$$g \leq a \text{hcod}$$

$$x = f \text{hcod} \wedge g \text{vcod},$$

then the following middle-four law about vertical and horizontal corestrictions holds:

$$([a|_*f]|_*[g|_*x]) = [(a|_*g)|_*(f|_*x)]$$

$$([a|_*f]|_*[g|_*x]) = [(a|_*g)|_*(f|_*x)]$$

$$\begin{array}{ccc}
 \begin{array}{ccc}
 \xrightarrow{\quad} & & \xrightarrow{\quad} \\
 \downarrow & a & \downarrow \\
 \xrightarrow{\quad} & & \xrightarrow{\quad}
 \end{array} & \geq & \begin{array}{c} \downarrow g \\ \text{gvcod} \end{array} \\
 \vee \wr & & \vee \wr \\
 \xrightarrow{f} f\text{hcod} & \simeq & x = f\text{hcod} \wedge \text{gvcod}
 \end{array}$$

Similarly,

$$(a) \quad [(a|_*g)|_*(f|_*x)] = ([a|_*f]|_*[g|_*x]).$$

$$(b) \quad ([x_*|g]_*|[f_*|a]) = [(x_*|f)_*|(g_*|a)].$$

$$(c) \quad [(x_*|f)_*|(g_*|a)] = ([x_*|g]_*|[f_*|a]).$$

## Definition (cont'd)

If  $e, f, g$  and  $h$  are objects, the following law about meets satisfying middle-four holds:

$$(e \wedge_h f) \wedge_v (g \wedge_h h) = (e \wedge_v g) \wedge_h (f \wedge_v h).$$

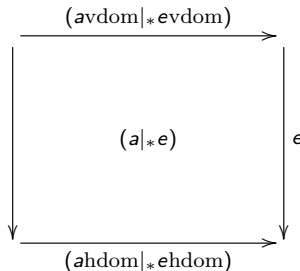
## Definition (cont'd)

If  $a$  is a double cell,  $e$  a vertical arrow and  $e'$  a horizontal arrow, then the following laws about domains and codomains preserving restrictions and corestrictions hold:

- (a)  $(a|_*e)\text{vdom} = (a\text{vdom}|_*e\text{vdom})$ .
- (b)  $(a|_*e)\text{vcod} = (a\text{vcod}|_*e\text{vcod})$ .
- (c)  $(e_*|a)\text{vdom} = (e\text{vdom}_*|a\text{vdom})$ .
- (d)  $(e_*|a)\text{vcod} = (e\text{vcod}_*|a\text{vcod})$ .
- (e)  $[a|_*e']\text{hdom} = [a\text{hdom}|_*e'\text{hdom}]$ .
- (f)  $[a|_*e']\text{hcod} = [a\text{hcod}|_*e'\text{hcod}]$ .
- (g)  $[e'_*|a]\text{hdom} = [e'\text{vdom}_*|a\text{hdom}]$ .
- (h)  $[e'_*|a]\text{hcod} = [e'\text{hcod}_*|a\text{hcod}]$ .

$$(a|_*e)\text{vdom} = (a\text{vdom}|_*e\text{vdom})$$

$$(a|_*e)\text{hdom} = (a\text{hdom}|_*e\text{hdom})$$



# Double Inductive Groupoids from Double Inverse Semigroups

## Construction (DIG)

*Given a double inverse semigroup  $(S, \odot, \odot)$ , we construct a double inductive groupoid*

$$\mathbf{DIG}(S) = (\mathbf{DIG}(S)_0, \text{Ver}(\mathbf{DIG}(S)), \text{Hor}(\mathbf{DIG}(S)), \text{Dbl}(\mathbf{DIG}(S)))$$

*as follows:*

**Objects:**  $\mathbf{DIG}(S)_0 = E(S, \odot) \cap E(S, \odot)$ .

## Construction (DIG cont'd)

**Vertical arrows:**  $\text{Ver}(\mathbf{DIG}(S)) = E(S, \odot)$ . Let  $u$  and  $v$  be any two vertically composable arrows:

- $uv\text{dom} = u \odot u^\odot$
- $uv\text{cod} = u^\odot \odot u$
- *Vertical composition:*  $u \bullet v = u \odot v$

**Horizontal arrows:**  $\text{Hor}(\mathbf{DIG}(S)) = E(S, \odot)$ . Let  $f$  and  $g$  be any two horizontally composable arrows:

- $f\text{hdom} = f \odot f^\odot$
- $f\text{hcod} = f^\odot \odot f$
- *Horizontal composition:*  $f \circ g = f \odot g$

## Construction (DIG cont'd)

$\text{Dbl}(\mathbf{DIG}(S)) = S(\odot, \odot)$ . Let  $a, b$  be any two horizontally composable double cells.

### Horizontally:

- $\text{ahdom} = a \odot a^\odot$
- $\text{ahcod} = a^\odot \odot a$
- *Horizontal composition:*  $a \circ b = a \odot b$
- *Horizontal partial order:*  $a \leq b$  iff  $a = \text{id}_e \odot b$  for some vertical arrow  $e$
- *Horizontal meet of two vertical arrows  $e$  and  $f$ :*  $e \wedge_h f = e \odot f$
- *If we have a vertical arrow  $e \leq \text{ahcod}$ , define  $(a|_*e) = a \odot e$*
- *If  $e \leq \text{ahdom}$ , define  $(e_*|a) = e \odot a$ .*

## Construction (DIG cont'd)

$\text{Dbl}(\mathbf{DIG}(S)) = S(\odot, \odot)$ . Let  $a, b$  be any two vertically composable double cells.

### Vertically:

- $\text{avdom} = a \odot a^\odot$
- $\text{avcod} = a^\odot \odot a$
- Vertical composition:  $a \bullet b = a \odot b$
- Vertical partial order:  $a \lesssim b$  iff  $a = 1_e \odot b$  for some horizontal arrow  $e$
- Vertical meet of two horizontal arrows  $e$  and  $f$  :  $e \wedge_v f = e \odot f$
- If we have a horizontal arrow  $e \lesssim \text{avcod}$ , define  $[a|_*e] = a \odot e$
- If  $e \lesssim \text{avdom}$ , define  $[e_*|a] = e \odot a$

# Double Inductive Groupoids from Double Inverse Semigroups

## Theorem

*If  $S(\odot, \odot)$  is a double inverse semigroup, then  $\mathbf{DIG}(S)$ , as constructed in Construction DIG, is a double inductive groupoid.*

# Double Inverse Semigroups from Double Inductive Groupoids

## Construction (DIS)

*Given a double inductive groupoid*

$$\mathcal{G} = (\text{Obj}(\mathcal{G}), \text{Ver}(\mathcal{G}), \text{Hor}(\mathcal{G}), \text{Dbl}(\mathcal{G})),$$

*we construct a double inverse semigroup **DIS**( $\mathcal{G}$ ) = ( $S, \odot, \odot$ ) as follows:*

- Its elements are the double cells of  $\mathcal{G}$ ;  $S = \text{Dbl}(\mathcal{G})$ .*

## Construction (DIS cont'd)

- For any  $a, b \in S$ , define

$$a \odot b = (a|_* a\text{hcod} \wedge_h b\text{hdom}) \circ (a\text{hcod} \wedge_h b\text{hdom}_* | b)$$

- For any  $a, b \in S$ , define

$$a \odot b = [a|_* a\text{vcod} \wedge_v b\text{vdom}] \bullet [a\text{vcod} \wedge_v b\text{vdom}_* | b]$$

# Double Inverse Semigroups from Double Inductive Groupoids

## Theorem

*If  $\mathcal{G}$  is a double inductive groupoid, then  $\mathbf{DIS}(\mathcal{G})$ , as constructed in Construction DIS, is a double inverse semigroup.*

Most of the work in proving this is in checking that the middle-four interchange law is satisfied.

$$\begin{array}{c}
 \boxed{a} \quad \boxed{b} \\
 \hline
 \boxed{c} \quad \boxed{d}
 \end{array}
 =
 \begin{array}{c}
 \boxed{a} \quad \boxed{b} \\
 \hline
 \boxed{c} \quad \boxed{d}
 \end{array}
 =
 \begin{array}{c}
 \boxed{a} \quad \boxed{b} \quad \boxed{c} \quad \boxed{d} \\
 \hline
 \boxed{c} \quad \boxed{d}
 \end{array}$$

# An Isomorphism of Categories

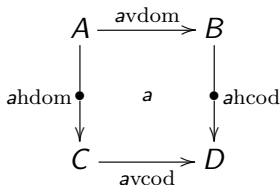
## Notation

*We denote the category of double inductive groupoids with double inductive functors as **DIG** and we denote the category of double inverse semigroups with double semigroup homomorphisms as **DIS**.*

## Theorem

*There exists an isomorphism of categories between **DIG** and **DIS**.*

Consider a double cell in a double inductive groupoid



Recall that domains and codomains may be written as semigroup products and that double inverse semigroups are commutative.

Then

- $a_h := ah\text{dom} = a \odot a^\odot = a^\odot \odot a = ah\text{cod}$
- $a_v := av\text{dom} = a \odot a^\odot = a^\odot \odot a = ah\text{cod}$

Similarly, the domain and codomain of a vertical or horizontal arrows are equal, so that

- $A = a_h \text{hdom} = a_h \text{hcod}$
- $A = a_v \text{vdom} = a_v \text{vcod}$

Ultimately,  $a$  is of the form

$$\begin{array}{ccc}
 A & \xrightarrow{a_v \text{dom}} & B \\
 \downarrow a_h \text{dom} & a & \downarrow a_h \text{cod} \\
 C & \xrightarrow{a_v \text{cod}} & D
 \end{array}
 =
 \begin{array}{ccc}
 A & \xrightarrow{a_h} & A \\
 \downarrow a_v & a & \downarrow a_v \\
 A & \xrightarrow{a_h} & A
 \end{array}$$

Let  $\mathcal{G}$  be a double inductive groupoid and let  $A$  be an object of  $\mathcal{G}$ .  
Then there is a natural collection of double cells

$$(A)\mathcal{S}_{\mathcal{G}} = \left\{ a \in \text{Dbl}(\mathcal{G}) \left| \begin{array}{ccc} A & \xrightarrow{a_h} & A \\ a_v \downarrow \bullet & a & \bullet \downarrow a_v \\ A & \xrightarrow{a_h} & A \end{array} \right. \right\}$$

- Recall: Objects of double inductive groupoids are idempotent with respect to both operations of its corresponding double inverse semigroup.
- Double inverse semigroups are commutative.

•

$$(a \odot b) \odot (a \odot b) = (a \odot a) \odot (b \odot b) = a \odot b$$

•

$$\begin{aligned}
 a \odot b &= (a \odot b) \odot (a \odot b) \\
 &= (a \odot b) \odot (b \odot a) \\
 &= (a \odot b) \odot (b \odot a) \\
 &= (a \odot b) \odot (a \odot b) \\
 &= a \odot b.
 \end{aligned}$$

## Lemma

*The vertical and horizontal order relations on the objects of a double inductive groupoid coincide.*

## Theorem

*These one-object double inductive groupoids are precisely Abelian groups.*

## Proposition

*Let  $\mathcal{G}$  be a double inductive groupoid. If  $A$  and  $B$  are objects in  $\mathcal{G}$  with  $A \leq B$ , then there is an Abelian group homomorphism*

$$\varphi_{A \leq B} : (B)\mathcal{S}_{\mathcal{G}} \rightarrow (A)\mathcal{S}_{\mathcal{G}}.$$

This discussion results in an **Ab**–valued presheaf

$$\mathcal{S}_{\mathcal{G}} : \text{Obj}(\mathcal{G})^{\text{op}} \rightarrow \mathbf{Ab}.$$

### Theorem

*Arbitrary double inverse semigroups are **Ab**–valued presheaves over meet-semilattices.*

## Construction

If  $P : L^{\text{op}} \rightarrow \mathbf{Ab}$  is a presheaf of Abelian groups on a meet-semilattice, define a double inductive groupoid  $\mathcal{G} = PF'$  with the following data:

**Objects:**  $\text{Obj}(\mathcal{G}) = L$

**Vertical/horizontal arrows:**

$\text{Ver}(\mathcal{G}) = \text{Hor}(\mathcal{G}) = \{e_A : A \rightarrow A : A \in L\},$

- $e_A$  is the group unit of the Abelian group  $AP$  for each  $A$  in  $L$ .
- (Co)domains:  $e_A \text{dom} = e_A \text{cod} = A$
- Composition:  $e_A \circ e_A = e_A \bullet e_A = e_A$ .
- Meets:  $e_A \wedge e_B = A \wedge B$  to be that from  $L$ .

## Construction (cont'd)

**Double cells:**  $\text{Db}(\mathcal{G}) = \coprod_{A \in L} AP$

- *Disjoint union of all Abelian groups  $AP$  for  $A$  in  $L$ .*
- *A double cell  $a$  is contained in an Abelian group  $AP$  for some  $A \in L$ .*
- $\text{ahdom} = \text{ahcod} = \text{avdom} = \text{vcod} = e_A$ .
- *Composites: group products*
- *If  $e_u \leq e_A = \text{ahdom}$ ,*
  - *Restriction of  $a$  to  $e_u$ :*

$$(e_u * | a) = e_u * _u (a) \varphi_{u \leq A} = (a) \varphi_{u \leq A}$$

- *Corestrictions are similarly defined.*

## Notation

*Denote the category of presheaves of Abelian groups on meet-semilattices by **AbMeetSLatt**.*

## Theorem

*The categories **DIG** and **AbMeetSLatt** are isomorphic.*

Recall:

- Kock showed double inverse semigroups are commutative.
- Double inverse semigroups are exactly presheaves of Abelian groups on meet-semilattices.

## Theorem

*Double inverse semigroups are commutative and improper. That is,  $(S, \odot, \oslash)$  is a double inverse semigroup if and only if both  $\odot$  and  $\oslash$  are commutative inverse semigroup operations with  $\oslash = \odot$ .*